

Geo (λ)/ Geo (μ) +G/2 Queues with Heterogeneous Servers Operating under FCFS Queue Discipline

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Received July 02, 2014; Revised March 01, 2015; Accepted March 06, 2015

Abstract This article discusses the steady analysis of a discrete time queue of Geo/Geo+G/2 type. All arriving customers are served either by server-1 according to a geometrically distributed service time $S_1=k$ slots for $k=1,2, \dots, \infty$, with mass function $f_1(k)=\Pr(S_1=k) = \mu(1-\mu)^{k-1}$ with mean rate $0 < \mu < 1$ or by server-2 with a general service time $S_2= k$ for $k=1,2, \dots, \infty$, with mass function $f_2(k)=\Pr(S_2=k)$ with mean service time is $\beta = \sum_{k=1}^{\infty} k f_2(k)$ or

mean service rate $\mu_2=1/\beta$. Sequel to some objections raised on the use of the classical 'First Come First Served (FCFS)' queue discipline when the two heterogeneous servers operate as parallel service providers, an alternative queue discipline in a serial configuration of servers are considered in this work; the objective is that if, in a single-channel queue in equilibrium, the service rate suddenly increases and exceeds the present service capacity, install a new channel to work serially with the first channel as suggested by Krishnamoorthy (1968). Using the embedded method subject to different service time distributions we present an exact analysis for finding the 'Probability generating Function (PGF)' of steady state number of customers in the system and most importantly, the actual waiting time expectation of customers in the system. This work shows that one can obtain all stationary probabilities and other vital measures for this queue under certain additional and simple but realistic assumptions.

Keywords: poisson arrival, service time distribution, pgf of queue length distribution, waittime distribution, mean queue length and mean waiting time

Cite This Article: Thaga Keaogile, Adebayo Fatai Adewole, and Sivasamy Ramasamy, "Geo (λ)/ Geo (μ) +G/2 Queues with Heterogeneous Servers Operating under FCFS Queue Discipline." *American Journal of Applied Mathematics and Statistics*, vol. 3, no. 2 (2015): 54-58. doi: 10.12691/ajams-3-2-2.

1. Introduction

Asynchronous transfer mode (ATM) multiplexers and broadband integrated services digital network (B-ISDN) use to transfer data sets, voice and video communications on a discrete time basis. Hence for studying the important characteristics of such discrete-time queueing of jobs served by a computer or a telecommunication device, the time axis can be divided into slots (fixed-length of contiguous intervals called slots of unit length (=right-end boundary-left-end boundary)).

Over the recent years, several authors Singh (1968), Hoksad (1979), Boxma et al. (2002), Efrosinin (2008), Kim et al. (2011), Krishnamoorthy and Srineevasan (2012) have studied continuous time queue length processes of either two server or multi-servers service systems. As the discrete-time queueing systems has been used to model computer and communication systems, authors Takagi (1993), and Bruneel and Kim (1993) have presented most of the basic features of discrete-time parallel queueing systems to that of continuous counterparts.

This paper discusses a special case of at the most only one arrival occurring (early arrival or late arrival) in a slot

at slot boundaries and service to a job that can only start at a slot boundary. The service duration of a job and the inter-arrival durations between consecutive job arrivals are measured as random number of slot durations. The proposed discrete-time queueing system is Geo(λ)/Geo(μ)+G/2 that has an infinite number of waiting positions with one faster server *i.e.* server-1 and a slow server *i.e.* server-2. One special feature of this investigation is that it derives those parallel results that have already been obtained to the corresponding continuous time version M/M+G/2 queues by Sivasamy and Kgosi (2014).

Late Arrival Model: Arrivals that come late in the slot (*i.e.* just before the slot boundary and before the service completions due to occur at the end of that slot) are called late arrivals. Number left behind in the queue as seen by a departure will include each late arrival and the time spent waiting in queue equals number of slots spent waiting for service not including the slot in which the job arrives.

Early Arrival Model: Arrivals that come early in a slot just after the left end slot boundary such that service completion of this arrival can occur just before the right-end slot boundary of the same slot. Number left behind in the queue as seen by a departure will not include the arrivals as in the late arrival cases. However the time spent

waiting in queue equals number of slots spent waiting for service including the slot in which the job arrives.

Geometric (or Bernoulli) Arrival Process: Assume that the numbers of jobs that arrive in successive slots are independent, identically distributed (i.i.d.) random variables subject to a condition that only one job can arrive in a slot with probability λ ($0 < \lambda < 1$) and that no jobs arrive in a slot with probability $1 - \lambda$.

It ensures that the inter-arrival time A is geometrically distributed with mean $1/\lambda$ and with probability the probability distribution $P\{A = k \text{ slots}\} = \lambda (1 - \lambda)^{k-1}$ for $k = 1, 2, \dots, \infty$ while $A(z)$ = probability generating function (PGF) $A(z)$ of inter-arrival times and $L(z)$ = the PGF of the number of arrivals in a slot are respectively given by

$$A(z) = \frac{\lambda z}{1 - (1 - \lambda)z}; \quad |z| < 1 \text{ and } L(z) = 1 - (1 - z)\lambda \quad (1)$$

Service time distributions: For customers serviced by server-1, their service time S_1 follows geometric distribution with mass function $f_1(k) = P(S_1 = k) = \mu (1 - \mu)^{k-1}$ for $k = 1, 2, \dots, \infty$ with mean $1/\mu$ or service rate μ ($0 < \mu < 1$) per slot and hence the PGF of S_1 and $L_s(z)$ = the PGF of the number of departures in a slot are respectively

$$F_1(z) = \sum_{k=1}^{\infty} \mu (1 - \mu)^{k-1} z^k = \frac{\mu z}{1 - (1 - \mu)z}; \quad (2)$$

$$|z| < 1 \text{ and } L_s(z) = 1 - (1 - z)\mu$$

Let the service times S_2 of customers serviced by server-2 follow a general distribution $f_2(k) = P(S_2 = k)$ with PGF is

$$F_2(z) = \sum_{k=1}^{\infty} f_2(k) z^k; \quad |z| < 1 \quad (3)$$

Both of these service times S_1 and S_2 are assumed to be mutually independent and each is independent of inter-arrival time distribution also. One of simple ways of connecting the servers in series subject to servicing of customers according to the FCFS queue discipline is proposed below:

- If an arriving job or customer enters into the idle system, his service is immediately initiated by server-1. This customer is then served by the server-1 at a constant rate μ if no other customer arrives during the on-going service period;
- Otherwise i.e. if at least one more customer arrives before the on-going service is completed then the same customer is served jointly by both servers according to the service time distribution $f_{\min}(k) = P(S_{\min} \leq k)$, where $S_{\min} = \text{Min}(S_1, S_2)$ and the PGF of S_{\min} is $F_{\min}(z)$:

$$\begin{aligned} F_{\min}(z) &= \sum_{k=1}^{\infty} \left[\frac{P(S_1 = k)P(S_2 \geq k)}{+P(S_1 > k)P(S_2 = k)} \right] z^k \\ &= \sum_{k=1}^{\infty} \left[\frac{P(S_1 = k) z^k - P(S_1 = k)P(S_2 < k)}{+P(S_1 > k)P(S_2 = k)} \right] z^k \\ &= F_1(z) + \sum_{k=1}^{\infty} P(S_2 = k) ((1 - \mu)z)^k \\ &\quad - \sum_{k=1}^{\infty} \mu (1 - \mu)^{k-1} \sum_{j=1}^{k-1} f_2(j) z^k \text{ since } P(S_1 > k) \end{aligned}$$

$$\begin{aligned} &= (1 - \mu)^k \\ &= F_1(z) + \sum_{k=1}^{\infty} P(S_2 = k) ((1 - \mu)z)^k \\ &\quad - (\mu z) \sum_{k=1}^{\infty} \{(1 - \mu)z\}^{k-1} \sum_{j=1}^{k-1} f_2(j) \\ &= F_1(z) + \sum_{k=1}^{\infty} P(S_2 = k) ((1 - \mu)z)^k \\ &\quad - (\mu z) \sum_{j=1}^{\infty} f_2(j) \sum_{k-1=j}^{\infty} \{(1 - \mu)z\}^{k-1} \\ &= F_1(z) + \sum_{k=1}^{\infty} P(S_2 = k) ((1 - \mu)z)^k \\ &\quad - \left(\frac{\mu z}{1 - (1 - \mu)z} \right) \sum_{j=1}^{\infty} f_2(j) \{(1 - \mu)z\}^j \\ &= F_1(z) + F_2((1 - \mu)z) \{1 - F_1(z)\} \end{aligned} \quad (4)$$

The expected (i.e. mean) values of S_1 , S_2 and of S_{\min} variables are respectively $E(S_1) = F_1'(z) = 1/\mu$, $E(S_2) = F_2'(1)$ and $E(S_{\min}) = F_{\min}'(1)$. For example, if the service time S_2 is a geometric random variable with mean $(1/\mu_2)$ then the mean of S_{\min} can be obtained from (4) i.e.

$$\begin{aligned} E(S_{\min}) &= F_{\min}'(1) = F_1'(1) \{1 - F_2(1 - \mu)\} \\ &= 1 / (\mu + \mu_2 - \mu\mu_2) \end{aligned}$$

It is then evident that the service rate of S_{\min} is $\mu_1 + \mu_2 - \mu_1\mu_2$ due to the fact that there is a positive probability of serving a job simultaneously by both servers just before the end boundary of a slot.

Lemma: If no job is served simultaneously by both servers with service time = k slots, the service (or departure) rate ' $r(k)$ ' of jobs offered by the combined service time distributions $f_1(k)$ and $f_2(k)$ of serialized servers of the $\text{Geo}(\lambda) / \text{Geo}(\mu) + G/2$ queueing system is given by

$$r(k) = \frac{f_1(k)}{1 - P(S_1 < k)} + \frac{f_2(k)}{1 - P(S_2 < k)} = \mu + \frac{f_2(k)}{1 - P(S_2 < k)}$$

Let the mean service times of server-2 and that of server-1 be $\beta = F_2'(1) = \sum_{k=1}^{\infty} k f_2(k) < \infty$ and

$$F_1'(1) = \sum_{k=1}^{\infty} k f_1(k) = \frac{1}{\mu} \quad (0 < \mu < 1) \text{ respectively. Thus the}$$

mean service rate of server-2 is $\mu_2 = 1/\beta$. Let $\rho = \frac{\lambda}{\mu} < 1$ be

used in the sequel. It is remarked that closed form expression to $r(k)$ can be obtained for particular cases of $f_2(k)$ like Geometric, Negative-binomial, or Phase (PH) type distributions. Assume that the service time distribution of the server-2 is one of these distributions with finite mean for our discussions to follow.

Negative-binomial Distribution: Let the service time S_2 of server-2 be Negative-binomial $\text{NB}(\alpha, \beta)$ random variable with mass function $b(k; \alpha, \mu_2) = P(S_2 = k)$. If S_2 denotes the number of slots required to complete the service by server-2 at the α^{th} success in a sequence

independent Bernoulli trials with probability of success β ($0 < \beta < 1$), then

$$b(k; \alpha, \beta) = \binom{k-1}{\alpha-1} \beta^\alpha (1-\beta)^{k-\alpha}; k = \alpha, \alpha+1, \dots, \infty$$

The mean of S_2 is $E(S_2) = (\alpha/\beta)$ and variance $V(S_2) = \alpha(1-\beta)/\beta^2$. The mean service rate $\mu_2 = \beta/\alpha$, a constant for a fixed and known α and β values. Then the service (or departure) rate 'r(k)' of jobs offered by the combined service time distributions $f_1(k)$ and $f_2(k)$ of serialized servers of the $Geo(\lambda)/Geo(\mu)+NB(\alpha, \beta)/2$ queueing system is given by

$$r(k) = \frac{f_1(k)}{1-P(S_1 < k)} + \frac{f_2(k)}{1-P(S_2 < k)}$$

$$= \mu + \mu_2 = \mu + \frac{\beta}{\alpha}$$

for all $k = \alpha, \alpha+1, \dots, \infty$

If the mean service time of server-2 i.e. $(1/\mu_2) \rightarrow \infty$ (tends to ∞) then his service rate $\mu_2 \rightarrow 0$ and thus $Geo(\lambda)/Geo(\mu)+G/2$ system $\rightarrow Geo(\lambda)/Geo(\mu)/1$ system.

Queueing models operating under the above (1) through (4) conditions are a discrete time class of $Geo(\lambda)/Geo(\mu)+G/2$ queues. We are motivated by the numerous physical applications and contributions of the above $Geo/Geo+G/2$ model since it can conceptualize well in the analysis of queues found in modern telecommunications, computer business centres, banking systems and similar service operations where human nature is evident.

Primary objective part of this study is to highlight the needs for designing serially connected system with two heterogeneous servers/machines which does not violate the FCFS principle in place of paralleled heterogeneous servers who violate the classical FCFS queue discipline. For instance, if there are two parallel clerks in a reservation counter who provide service with varying speeds then customers might prefer to choose the fastest service provider among the free servers. On the other hand if the customer chooses the slowest server among the free servers then the one who could come subsequently may clear out earlier by obtaining service from the server providing service with faster rates what is exactly called a violation of the FCFS discipline. Hence there is a need for the designing of alternative ways to the parallel service providers that would reduce the impacts of violating the FCFS so that the resulting waiting times of customers are identical.

The PGFs of the number of customers present in the system and the waiting time distribution and their mean values have been obtained in section 2. A numerical illustration is also provided to support the results on mean waiting times. Section 3 highlights the various special features of the proposed methodology and its future scope.

2. Steady State Characteristics of $Geo(\lambda)/Geo(\mu)+G/2$ Queues with Service Time Distribution $F_{min}(K)$

We consider here 'the embedded time points' generated at the departure instants of jobs just after a service completion either by server-1 or by server-2. Hence the sequence of system states observed at these embedded points where the state of the system is represented by the number, $N_k = N(t_k)$, of jobs left behind in the queue by the k^{th} departing job at the departure epoch ' t_k ' forms a Markov Chain process $\{N_k\}$ on the state space $\mathbb{S} = \{0, 1, \dots, \infty\}$.

2.1. PGF of $\{N_k\}$

Let q_j be the steady state probability of finding 'j' customers in the system as observed by a departing customer and the z-transform of the probability

distribution of $\{q_j; j=0, 1, \dots, \infty\}$ be $V(z) = \sum_{j=0}^{\infty} q_j z^j$.

Let α_j denote the probability of 'j' arrivals in a service completion period with mass function $f_{min}(k)$ and δ_j denote the probability of 'j' arrivals in the geometrically distributed service time mass function $f_1(k)$. Since the arrivals come from a Geometric process at rate λ , we get, for $j=1, 2, 3, \dots, \infty$ that

$$\alpha_j = \sum_{k=1}^{\infty} f_{min}(k) \binom{k}{j} \lambda^j (1-\lambda)^{k-j} \tag{5}$$

and

$$\delta_j = \sum_{k=1}^{\infty} f_1(k) \binom{k}{j} \lambda^j (1-\lambda)^{k-j} \tag{6}$$

Let the respective z-transforms of the probability distributions $\{\alpha_j\}$ and $\{\delta_j\}$ be $A_{min}(z)$ and $A_1(z)$:

$$A_{min}(z) = \sum_{j=0}^{\infty} \alpha_j z^j = F_{min}(1-\lambda+\lambda z) \tag{7}$$

and

$$A_1(z) = \sum_{j=0}^{\infty} \delta_j z^j = F_1(1-\lambda+\lambda z) \tag{8}$$

Focusing on the embedded points under equilibrium conditions, let the unit step conditional transition probability of the system going from state 'i' of the $(k-1)^{st}$ embedded point to state 'j' in the k^{th} embedded point be $q_{ij} = P(N_k = j | N_{k-1} = i)$ for $i, j \in \mathbb{S}$. These transition probabilities will form the unit step transition probability matrix $Q = (q_{ij})$ as below:

$$Q = \begin{pmatrix} \delta_0 & \delta_1 & \delta_2 & \delta_3 & \dots & \dots & \dots & \dots \\ \delta_0 & \delta_1 & \delta_2 & \delta_3 & \dots & \dots & \dots & \dots \\ 0 & \alpha_0 & \alpha_1 & \alpha_2 & \alpha_3 & \dots & \dots & \dots \\ 0 & 0 & \alpha_0 & \alpha_1 & \alpha_2 & \alpha_3 & \dots & \dots \\ 0 & 0 & 0 & \alpha_0 & \alpha_1 & \alpha_2 & \alpha_3 & \dots \end{pmatrix} \tag{9}$$

Thus equilibrium state probabilities at the departure instants are given by $q_j = \lim_{n \rightarrow \infty} q_{ij}^n$ where q_{ij}^n represents the n-step probability of moving from state 'i' to 'j'. Let $q = (q_0, q_1, q_2, \dots)$ be a row vector and let $e = (1, 1, \dots, 1)$ be a column vector of unit elements of infinite order. Assume that $A_{min}'(1) < 1$, then the stationary distribution of the state

transition matrix Q exists and is given by the unique solution of the following system of equations:

$$q Q = q \text{ and } q e = 1 \tag{10}$$

Multiplying the j^{th} ($j=0, 1, 2, \dots$) equation of $qQ=q$, of (10) by z^j and summing all the left-hand sides and the right-hand sides from $j=0$ to $j=\infty$, we get the PGF (Probability generating Function) $V_{\min}(z)$ of the queue length distribution $\{q_j\}$ of the sequence $\{N_k\}$.

$$V_{\min}(z) = \frac{q_1 z [A_{\min}(z) - A_1(z)] + q_0 [A_{\min}(z) - A_1(z)z]}{A_{\min}(z) - z} \tag{11}$$

Since $q_1 = \rho q_0$, $A_{\min}'(1) < 1$, $A_1'(1) = \rho$ and $V(1) = 1$, we derive from (11) that

$$q_0 = \frac{1 - A'_{\min}(1)}{1 + (\rho - A'_{\min}(1))(1 + \rho)} \tag{12}$$

Thus

$$V_{\min}(z) = \frac{(z-1)A_1(z) + (1+\rho z)[A_1(z) - A_{\min}(z)]}{z - A_{\min}(z)} \tag{13}$$

$$\frac{1 - A'_{\min}(1)}{1 + (\rho - A'_{\min}(1))(1 + \rho)}$$

The mean number $E(N)$ of customers present in the system at a random point or at a departure epoch of time is

$$E(N) = \frac{[\rho - A'_{\min}(1)][\rho + A'_{\min}(1)(1 + \rho)]}{1 + \{\rho - A'_{\min}(1)\}(1 + \rho)} \tag{14}$$

$$+ \rho + \frac{A'_{\min}(1)^2}{1 - A'_{\min}(1)}$$

The discrete time equivalent of the PASTA (Poisson Arrivals See Time Averages) is referred to as BASTA (Bernoulli Arrivals See Time Averages) or GASTA (Geometric Arrivals See Time Averages). Using this BASTA property, it can be claimed that the distribution given by $V_{\min}(z)$ of (13) will also hold for the number of jobs found in the system as seen by an arriving job.

Denote the PGF of the waiting time W that equals the number of slots for which an arriving job stays in the system with its distribution $W(k) = P(W = k)$ by $W(z)$ for $0 < z < 1$. Then replacing $1 - (\lambda - \lambda z)$ by z in (13) one can show that

$$W(1 - \lambda + \lambda z) = V_{\min}(z) \tag{15}$$

The mean waiting time \bar{W} of a customer in the system obtained from (15) is found to satisfy the well-known Little's formula $\lambda \bar{W} = E(N)$.

2.2. How $Geo(\lambda) / Geo(\mu) + G/2 \rightarrow Geo(\lambda) / Geo(\mu) / 1$

For the purpose of proving the behaviour $Geo(\lambda) / Geo(\mu) + G/2 \rightarrow Geo(\lambda) / Geo(\mu) / 1$, the PGFs of queue length and the waiting time distributions of $Geo(\lambda) / Geo(\mu) / 1$ system, assuming the load $\rho = (\lambda/\mu) < 1$, are respectively given by $Q_0(z)$ and $W_0(z)$ below:

$$Q_0(z) = \frac{(1-z)(1-\rho)F_1(1-\lambda+\lambda z)}{F_1(1-\lambda+\lambda z)-z}; |z| < 1$$

with mean

$$E(Q_0) = \rho + \frac{\lambda^2}{2(1-\rho)} E(S_1(S_1-1))$$

$$W_0(1-\mu+\mu z) = Q_0(z) \text{ or}$$

$$W_0(z) = \frac{(1-z)(1-\rho)F_1(z)}{(1-z)-\lambda(1-F_1(z))}; |z| < 1$$

with mean

$$\bar{W}_0 = E(Q_0) / \lambda$$

Lemma 2: If the mean service time $(1/\mu_2)$ of the general distribution G of the second server does not exist i.e. $(1/\mu_2) \rightarrow \infty$ as $\mu_2 \rightarrow 0$, it is shown that

$$Geo(\lambda) / Geo(\mu) + G / 2 Geo(\lambda) / Geo(\mu) / 1$$

Proof: If $(1/\mu_2) \rightarrow \infty$ as $\mu_2 \rightarrow 0$ then $F_2(z)$ cannot exist and thus which is replaced by a zero in $F_{\min}(z) = F_1(z) + F_2((1-\mu)z)\{1-F_1(z)\}$; this in turn proves that $A_{\min}(z) = F_1(1-\lambda+\lambda z) = A_1(z)$. Proof of the lemma follows from replacing $A_{\min}(z)$ by $A_1(z)$ in the PGF of the queue length distribution $V_{\min}(z)$ of (13) of the $Geo(\lambda) / Geo(\mu) + G/2$ since $V_{\min}(z) = \frac{(z-1)(1-\rho)A_1(z)}{z - A_1(z)}$

3. Shortest Delay Time Environment For $Geo(\lambda) / Geo(\mu) + G/2$ queues

The forgoing embedded methodology of analysing the embedded queue length process of the $\{N_k\}$ sequence of $Geo(\lambda) / Geo(\mu) + G/2$ queues through the PGF $V_{\min}(z)$ queues can also be extended to design a shortest processing environment as outlined below.

Let $f_{\max}(k) = P(S_{\max} < k)$, where S_1 and S_2 denote the service times of server-1 and server-2 respectively as already defined and $S_{\max} = \text{Max}(S_1, S_2)$ and the PGF of T_{\max} be $F_{\max}(z)$. Then

$$F_{\max}(z) = \sum_{k=1}^{\infty} f_{\max}(k) z^k = F_2(z) - F_2(z(1-\mu))\{1-F_1(z)\}$$

Considered below in (16) is the new joint service time PGF of 'f(k)', the probability mass function of the joint service time of the two independent servers who are serially connected to serve the jobs without violating the FCFS discipline

$$f(k) = \pi_1 f_{\min}(k) + \pi_2 f_{\max}(k) \tag{16}$$

where π_1 and π_2 are the probability values of choosing the service time PDFs $f_{\min}(k)$ and $f_{\max}(k)$ respectively. The corresponding z-transform of number of arrivals during the joint service PGF of f(k) is

$$A(z) = \sum_{j=0}^{\infty} \alpha_j z^j = \pi_1 F_{\min}(1-\lambda+\lambda z) \tag{17}$$

$$+ \pi_2 F_{\max}(1-\lambda+\lambda z)$$

Let $\rho_2 = \frac{\lambda}{\mu_2}$ and $\rho = \frac{\lambda}{\mu}$. Assume that $A'(1) < 1$ of (17),

i.e.

$$A'(1) = \pi_1 \rho \{1 - F_2(1 - \mu)\} + \pi_2 \{\rho_2 + \rho F_2(1 - \mu)\} < 1 \quad (18)$$

Now on substituting $A(z)$ of (17) in place of $A_{\min}(z)$ of (13), one can verify that $V(z)$, the PGF of the queue length distribution of the sequence $\{N_k\}$ ($V(z)$ replaces $V_{\min}(z)$ of (13) and $f(k)$ in place of $f_{\min}(k)$) is

$$V(z) = \frac{(z-1)A_1(z) + (1+\rho z)[A_1(z) - A(z)]}{z - A(z)} \quad (19)$$

$$\frac{1 - A'(1)}{1 + (\rho - A'(1))(1 + \rho)}$$

Lemma: If $\lambda \neq \mu \neq \mu_2$ and $\pi_1 < \pi_2$ then $A'(1) < 1$ of (17)

implies that $\rho + \rho_2 < \frac{1}{\pi_1}$.

Proof: Notice that $F_2(1-\mu) = \mu_2(1-\mu)/(1-(1-\mu)(1-\mu_2)) = \mu/(\mu + \mu_2 - \mu\mu_2)$ and

$$A'(1) = \pi_1 \rho \left\{ \frac{\mu_2(1-\mu)}{\mu + \mu_2 - \mu\mu_2} \right\} + \pi_2 \left\{ \rho_2 + \rho \frac{\mu}{\mu + \mu_2 - \mu\mu_2} \right\}$$

If $A_1'(1) < 1$, then

$$\pi_1 \rho \left\{ \frac{\mu_2(1-\mu)}{\mu + \mu_2 - \mu\mu_2} \right\} + \pi_1 \left\{ \rho_2 + \rho \frac{\mu}{\mu + \mu_2 - \mu\mu_2} \right\} < 1 \quad ;$$

hence the lemma.

Conditional Average Waiting Time(s): Let w_1 and w_2 be the mean waiting times corresponding to $\mu < \mu_2$ and $\mu > \mu_2$ values respectively. Thus $w_1 = E(N)/\lambda$ when $\mu < \mu_2$ and $w_2 = E(N)/\lambda$ when $\mu > \mu_2$. Now $E(N) = V'(1)$ from (19) is

$$E(N) = \rho + \frac{A'(1)^2}{1 - A'(1)} + \frac{\{\rho - A'(1)\}[\rho + A'(1)(1 + \rho)]}{1 + \{\rho - A'(1)\}(1 + \rho)} \quad (19)$$

From the design criterion of the above model, it is expected that if the service rate μ of server-1 is larger than that μ_2 of server-2 then $w_2 < w_1$ with increasing values of π_1 . For an illustration, assuming that the probability of serving a job simultaneously by both servers in the same slot is zero (i.e. $\mu\mu_2 = 0$) w_1 is calculated fixing $\lambda = 0.48$, $\mu = 0.45$ and $\mu_2 = 0.8$ (geometric case) and w_2 is computed fixing $\lambda = 0.48$, $\mu = 0.8$ and $\mu_2 = 0.45$ (geometric case) while π_1 varies between 0.4 and 1 and these values are reported in Table-1 (which also ensures that $w_2 < w_1$ uniformly).

Table 1. Conditional Mean waiting times w_1 when $(\lambda=0.48, \mu=0.45, \mu_2=0.8)$ and w_2 when $(\lambda=0.48, \mu=0.8, \mu_2=0.45)$ of Geo(λ)/ Geo(μ)+Geo(μ_2)/2 queues

π_1	$\lambda=0.48, \mu=0.45, \mu_2=0.8$			$\lambda=0.48, \mu=0.8, \mu_2=0.45$		
	q_0	$E(N)$	w_1	q_0	$E(N)$	w_2
0.40	0.059237	12.49346	26.02805	0.159046	10.30738	21.47371
0.50	0.112444	5.672364	11.81742	0.265957	4.046809	8.430851
0.60	0.153802	3.725747	7.761973	0.332964	2.421429	5.044645
0.70	0.186872	2.838598	5.913746	0.378894	1.736749	3.618226
0.80	0.213920	2.345471	4.886397	0.412340	1.379927	2.874848
0.90	0.236452	2.038537	4.246951	0.437783	1.169337	2.436119
0.95	0.246364	1.926196	4.012908	0.448350	1.094933	2.281110

This is a kind of two-server queueing application subject to unequal service rates will lead to the event of shortest mean waiting time or delay only if the service rate μ of server-1 is larger than that μ_2 of server-2.

3.1. Conclusion

As the analysis of the above Geo(λ)/ Geo(μ)+G(μ_2)/2 queues seems to be simpler and it is an easy procedure to implement if any real life application warrants. We conclude that the proposed method so far discussed is a good alternative for use instead of fitting of Geo(λ)/ Geo(μ), G(μ_2)/2 parallel server queues. Most importantly there are few new results produced from above work that could be applied on the two-serially connected machines in the field of engineering problems rather than applied science areas. To extract more information like arrival epoch probabilities one can employ Markov Renewal theory as in Senthamarikkannan and Sivasamy [9] and [10].

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